Data-intensive computing systems



Hbase MongoDB

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Introduction

- □ HDFS focuses on write once, read many workloads
- □ What if we need to store data and, in addition to full scans for analytics, occasionally update and read random elements?
- □ Example: Webtable
 - Output of the web crawler: for each web page URL, store the page, the attributes (language, MIME types), ...
 - The whole table is accessed by MapReduce jobs for analytics
 - Random rows are updated by the crawler
 - Random rows are read by the crawler or by other systems

□ Solution: HBase

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HBase: Overview

- □ HBase is a distributed column-oriented data store built on top of HDFS
- HBase is an Apache open source project whose goal is to provide storage for the Hadoop Distributed Computing
- Data is logically organized into tables, rows and columns



HBase: Part of Hadoop's Ecosystem



HBase vs. HDFS

□ Both are distributed systems that scale to hundreds or thousands of nodes

□ HDFS is good for batch processing (scans over big files)

- Not good for record lookup
- Not good for incremental addition of small batches
- Not good for updates
- □ HBase is designed to efficiently address the above points
 - Fast record lookup
 - Support for record-level insertion
 - Support for updates (not in place)
- □ HBase updates are done by creating new versions of values



HBase vs. HDFS (Cont'd)

	Plain HDFS/MR	HBase
Write pattern	Append-only	Random write, bulk incremental
Read pattern	Full table scan, partition table scan	Random read, small range scan, or table scan
Hive (SQL) performance	Very good	4-5x slower
Structured storage	Do-it-yourself / TSV / SequenceFile / Avro / ?	Sparse column-family data model
Max data size	30+ PB	~1PB

If application has neither random reads or writes -> Stick to HDFS



HBase Data Model

- HBase is based on Google's Bigtable model
 - Key-Value pairs

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- Rows are ordered by keys
- □ Table schema only define it's column families
 - Each family consists of any number of columns
 - Each column consists of any number of versions
 - Columns only exist when inserted,
 - Everything except table names are byte[]
 - (Row, Family: Column, Timestamp) \rightarrow Value



HBase: Keys and Column Families



Notes on Data Model

- □ HBase schema consists of several Tables
- □ Each table consists of a set of Column Families
 - Once the Column Families are defined, Columns can be added at any time
 - HBase has Dynamic Columns
 - Hardcoded name convention
 - column_family_name:column_ID
- $\hfill\square$ Not all the columns are used by all the keys
 - Table can be very sparse, many cells are empty





HBase Physical Model

- □ Each column family is stored in a separate file (called HTables)
- □ Key & Version numbers are replicated with each column family
- Empty cells are not stored

HBase maintains a multilevel index on values: <key, column family, column name, timestamp>

Table 5.3. ColumnFamily contents

Row Key	Time Stamp	ColumnFamily "contents:"
"com.cnn.www"	t6	contents:html = " <html>"</html>
"com.cnn.www"	ť5	contents:html = " <html>"</html>
"com.cnn.www"	t3	contents:html = " <html>"</html>

Table 5.2. ColumnFamily anchor

Row Key	Time Stamp	Column Family anchor
"com.cnn.www"	t9	anchor:cnnsi.com = "CNN"
"com.cnn.www"	t8	anchor:my.look.ca = "CNN.com"

Example

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Row key	Data
cutting	≈info: { 'height': '9ft', 'state': 'CA' } ≈roles: { 'ASF': 'Director', 'Hadoop': 'Founder' }
tlipcon	info: { 'height': '5ft7, 'state': 'CA' } roles: { 'Hadoop': 'Committer'@ts=2010,

info Column Family

Row key	Column key	Timestamp	Cell value
cutting	info:height	1273516197868	9ft
cutting	info:state	1043871824184	CA
tlipcon	info:height	1273878447049	5ft7
tlipcon	info:state	1273616297446	CA

roles Column Family

_	Row key	Column key	Timestamp	Cell value
	cutting	roles:ASF	1273871823022	Director
on disk by	cutting	roles:Hadoop	1183746289103	Founder
Row key, Col	tlipcon	roles:Hadoop	1300062064923	PMC
descending	tlipcon	roles:Hadoop	1293388212294	Committer
timestamp	tlipcon	roles:Hive	1273616297446	Contributor

Milliseconds since unix epoch



cloudera

Column Families

- Different sets of columns may have different properties and access patterns
- □ Configurable by column family:
 - Compression (none, gzip, ...)
 - Version retention policies
- □ Column Families stored separately on disk
 - Access one without wasting I/O on the other



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HBase Regions

□ Each HTable (column family) is partitioned horizontally into regions

- Regions are counterpart to HDFS blocks

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Each will be one region



HBase Components

□ Region

- A subset of a table's rows, like horizontal range partitioning
- Automatically done
- □ RegionServer (many slaves)
 - Manages data regions
 - Serves data for reads and writes (using a log)

□ Master

- Responsible for coordinating the slaves
- Assigns regions, detects failures
- Admin functions

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Big Picture



Access to data

□ Set of APIs that can be used to do basic operations

□ Put()

- Insert a new record (with a new key), or insert a record for an existing key
- The column ID (given an existing column family) must be also specified
- Implicit or explicit versioning
- □ Get()
 - retrieve the value given a key and a column ID
- □ Scan()
 - retrieve all the values given a column identifier (and a range of keys)
- Delete()

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- Multiple levels
 - Can mark an entire column family as deleted
 - Can make all column families of a given row as deleted



Access to data: Get()

Select value from table where key= 'com.apache.www' AND label= 'anchor:apache.com'

Row key	Time Stamp	Column "anchor:"	
	t12		
"com.apache.www"	tl1		
	t10	"anchor:apache.com"	"APACHE"
"com.cnn.www"	t9	"anchor:ennsi.com"	"CNN"
	t8	"anchor:my.look.ca"	"CNN.com"
	t6		
	t5		
	t3		



Access to data: Scan()

Time Row key Column "anchor:" Stamp t12 t11 "com.apache.www" "anchor:apache.com" "APACHE" t10 "anchor:cnnsi.com" "CNN" t9 "anchor:my.look.ca" "CNN.com" t8 "com.cnn.www" t6 t5 t3

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HBase Deployment





HBase vs. RDBMS

	RDBMS	HBase
Data layout	Row-oriented	Column-family- oriented
Transactions	Multi-row ACID	Single row only
Query language	SQL	get/put/scan/etc *
Security	Authentication/Authorizatio n	Work in progress
Indexes	On arbitrary columns	Row-key only
Max data size	TBs	~1PB
Read/write throughput limits	1000s queries/second	Millions of queries/second
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MongoDB



Overview

- Document oriented, not table/row oriented
- □ Collection of binary JSON (BSON) documents
- Schemaless
- No relations or transactions native in database
- □ Scalable and high-performance
- □ Full index support
- □ Written in C++
- □ Servers for all major platforms
- Drivers for all major development environments
- □ Free and open-source, but also commercial support

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BSON

- □ Binary JSON
- D Binary encoded serialization of JSON-like documents
- □ Like JSON, BSON supports the embedding of documents and arrays within other documents and arrays. BSON also contains extensions that allow representation of data types that are not part of the JSON spec. For example, BSON has a Date type and a BinData type.
- □ The driver performs translation from the language's "domain object" data representation to BSON, and back



MondoDB Data Model and Types

🖵 Data Model

- A MongoDB deployment hosts a number of databases
- A database holds a set of collections
- A collection holds a set of documents
- A document is a *set* of key-value pairs

Data types

- Basic: Null, Boolean, Integer (32- and 64-bit), Floating point, String
- More complex: Date, Code (JavaScript), Array, Embedded document



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Sample Document





Core MongoDB Operations

□ CRUD: create, read, update, and delete

Insert

- One at a time: db.mycollection.insert(mydoc)
- Batch insert
- □ Querying
 - Use find()/findOne() functions and a query document
 - Ranges, set inclusion, inequalities using \$ conditionals
- Delete
 - Documents that match some predicate, e.g. to remove the document just added

```
db.mycollection.remove({"_id": 1})
```

All documents in a collection
 db.mycollection.remove()

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Find

```
Get the entire collection (called posts)
  db.posts.find()
Get the documents satisfying some constraints
  db.posts.find({"author": "Kevin"})
GSpecifying Which Keys to Return
  db.mydoc.find({}, {"author", "tags"})
  {
    __id: ObjectId("4c4ba5c0672c685e5e8aabf3"),
    author: "Kevin",
    tags: [ "tech", "databases"]
  }
```



Ranges, Negation, OR-clauses



Limits, Skips, Sort, Count

- Limits the number of results to 3
 db.posts.find().limit(3)
- Skips the first three results and returns the rest db.posts.find().skip(3)
- □ Sorts by author ascending (1) and title descending (-1) db.posts.find().sort({"author":1, "title": -1})
- Counts the number of documents in the people collection matching the find(...)

```
db.people.find(...).count()
```

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Summary of MongoDB

- □ MongoDB is an example of a document-oriented NoSQL solution
- □ The query language is limited, and oriented around "collection" (relation) at a time processing
 - Joins are done via a query language
- The power of the solution lies in the distributed, parallel nature of query processing
 - Replication and sharding

