

# MATHEMATICAL FOUNDATIONS OF REINFORCEMENT LEARNING

## LECTURE IV: APPROXIMATE DYNAMIC PROGRAMMING

SUMMARY:

- COMPUTING OPTIMAL FEEDBACKS ONLINE
- APPROXIMATE VALUE ITERATION
- APPROXIMATE POLICY ITERATION
- BELLMAN EQN. METHODS.

REFS:

- CHAPTER 5, BERTSEKAS BOOK ON RL.

- A. ALLA, M. FALCONE, J.K. "AN EFFICIENT POLICY ITERATION

ALGORITHM FOR DYNAMIC PROGRAMMING EQUATIONS", SISC, 2015.

## PREVIOUS LECTURES:

INFINITE HORIZON COST:

$$J^N(i) = \lim_{N \rightarrow \infty} E \left[ \sum_{k=0}^N \alpha^k g(i_k, u_k, i_{k+1}) \mid i_0 = i \right]$$

*discount*  
*run next*  
*cost*

VALUE FUNCTION :  $J^*(i) = \min_u J^N(i)$



V.F. SATISFIES BELLMAN EQUATION:

$$J^*(i) = \min_{u \in \mathcal{U}(i)}$$

$$\sum_{j \in X} p_{ij}(u) (g(i, u, j) + \alpha J^*(j))$$

$\stackrel{+1}{\nearrow}$   
 $\stackrel{0}{\nearrow}$   
 $\downarrow$

VALID  $\forall i \in X$

WITH BELLMAN OPERATOR

$$TJ := \min_{u \in \mathcal{U}(i)} \sum_{j \in X} \bar{p}_{ij}(u) (\bar{g}(i, u, j) + \alpha J(j)) \Rightarrow J^* = TJ^*$$

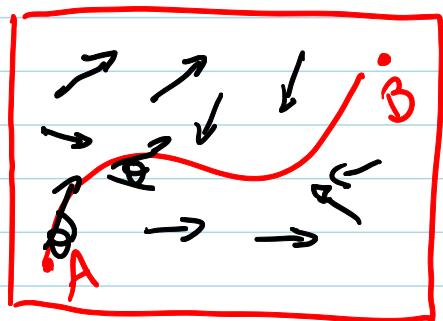
$T$  is a CONTRACTION Map  $\Rightarrow$  FIXED Point ITERATION

$$\underline{J}^*(i) = \lim_{K \rightarrow \infty} \underline{J}^*(i)$$
 WHERE  $\underline{J}^{K+1} = T \underline{J}^K$

Once  $\underline{J}^*$  has been computed,  $\mu^*$  is recovered as

$$\mu^*(i) \in \arg\min_{u \in U(i)} \sum_j p_{ij}(u) (q_i(u, j) + \alpha \underline{J}^*(j))$$

EXAMPLE : ZERNERO NAVIGATION PROBLEM



GO FROM A TO B IN MIN TIME SUBJECT TO WIND

DISCRETE EQNS OF MOTION

$$x_{k+1} = x_k + \Delta t (W_{\text{IND}}_x + \bar{V} \cos(\theta))$$

CONTROL  
 $(x_k, y_k)$

$$y_{k+1} = y_k + \Delta t (W_{\text{IND}}_y + \bar{V} \sin(\theta_k))$$

STATE  
 $(x_k, y_k)$

# SOME IDEAS FOR ANNs IN THE CONTROL FIELD

$$\mu^*(i) \in \arg\min \sum_j p(j)(u) (g(i, u, j) + \alpha \underbrace{J^*(j)}_{\text{REPLACE BY A LOW COMPLEXITY ANN}})$$

REPLACE BY A LOW COMPLEXITY ANN  
 $\tilde{J}(j, r)$

EXAMPLE: LQ Control

$$x_{k+1} = ax_k + bu_k, \quad x_k, u_k \in \mathbb{R}$$

$$g(x, u, y) = K u^2 + x^2, \quad K \text{ constant}$$

ASSUME A LINEAR ARCHITECTURE :=  $\tilde{J}(i, r) = r_0 + r_1 x + r_2 x^2$

$$\mu^*(x_i) \in \arg\min_u$$

$$K u^2 + x_i^2 + r_0 + r_1 (\underbrace{ax_i + bu}_{} + \underbrace{r_2 (ax_i + bu)^2}_{g(u)})$$

$$g(u) \Rightarrow g'(u) = 0$$

How do we incorporate the approx. of  $J^*(i)$  by  $\tilde{J}(i, \pi)$

↳ APPROXIMATE VALUE ITERATION:

$$\hat{J}_{k+1}(i) = \min_{\pi} \sum_j p_{ij}(u) (g(i, u, j) + \alpha \underbrace{\tilde{J}(j, \pi_k)}_{\text{red line}}),$$

$\forall i \in S_k$ , THEN  $R_{k+1}$  is UPDATED according to

$$\pi_{k+1} \in \arg \min_{\pi} \sum_{i \in S_k} \| \hat{J}_{k+1}(i) - \tilde{J}(i, \pi) \|^2$$

↳ DEFINE ARCHITECTURE  $\tilde{J}(i, \pi)$  ✓

↳ INITIAL PARAMETRIC GUESS  $\pi_0$  ✓

↳ UPDATE  $\hat{J}_{k+1}$  USING  $\tilde{J}(i, \pi)$  OVER A SUBSET  $S_k$

↳ USE  $\hat{J}_{k+1}(i)$  TO TRAIN  $\pi_{k+1}$

OBS: - THE ARCHITECTURE  $\tilde{J}(i, r)$  REMAINS, ONLY  $r$  CHANGES  
AFTER EVERY ITERATION.

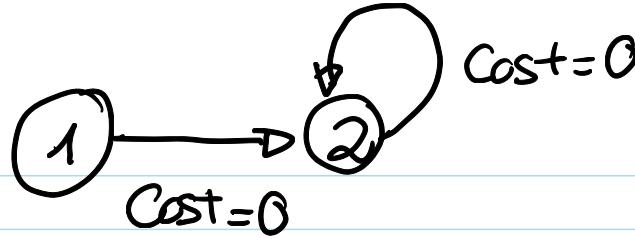
- THE EVALUATION OF  $P_{ij}$ 'S CAN BE COSTLY UNLESS #  $P_{ij} \neq 0$  IS SMALL.

OBS: AVI IS VERY SIMILAR TO VI, BUT IT IS CRITICAL TO ENSURE THAT

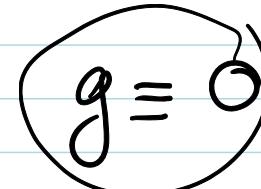
$$\boxed{\| \tilde{J}(i, \pi_{K+1}) - \tilde{J}_{K+1}(i) \| \leq \delta}$$

WITH  $\delta$  SUFFICIENTLY SMALL, OTHERWISE THE CONTRACTIVITY OF THE BELLMAN OPERATOR IS LOST

PATHOLOGICAL CASE:



$$\Rightarrow p_{11} = 0, p_{12} = 1 \\ p_{22} = 1, p_{21} = 0$$



Cost=0

$\alpha > 0$  discount factor

It's clear that  $J^*(1) = J^*(2) = 0$  (because  $g \geq 0$ )

USE BELLMAN EQN.

$$J^*(i) = \sum_j p_{ij} (0 + \alpha J^*(j))$$

BELLMAN

$$J^*(1) = \alpha J^*(2)$$

$$J^*(2) = \alpha J^*(2)$$

VALUE ITERATION

$$J_{k+1}(1) = \alpha J_k(2)$$

$$J_{k+1}(2) = \alpha J_k(2)$$

ASSUME A LINEAR APPROXIMATION  $\tilde{J}(i, r) = \underline{i} \cdot \underline{r}$

$$\Rightarrow J_k(1) = \pi_k, J_k(2) = 2 \cdot \pi_k.$$

$$\Rightarrow \overline{T} \tilde{J}_k = (2\alpha \pi_k, 2\alpha \pi_k) \Leftrightarrow \overline{\tilde{J}}_{k+1}$$

LINEAR LEAST SQUARES:

$$\pi_{k+1} \in \underset{\pi}{\operatorname{argmin}} \sum_{\tilde{J}(1, r)} \left( \underline{r} - 2\alpha \pi_k \right)^2 + \left( \underline{2r} - 2\alpha \pi_k \right)^2$$

$$\Rightarrow 2(r - 2\alpha \pi_k) + 2 \cdot (2r - 2\alpha \pi_k) \cdot 2 = 0$$

$$\pi_{k+1} = \frac{6}{5} \alpha \pi_k$$

$$\Rightarrow \text{WE NEED } \alpha < \frac{5}{6} \text{ TO CONVERGE } (\pi^* = 0)$$

# APPROXIMATE POLICY ITERATION

PI: Given  $\pi_k$ , solve

$$J^{k+1} = T_{\pi^k} J^{k+1} \quad (\text{Policy EVAL})$$

$$\pi^{k+1} = \arg \max_{\pi} \sum p_{ij}(u) (g(i, a_j) + J^{k+1}(j)) \quad (\text{Policy UPDATE})$$

↳ Incorporate ANN By replacing  $J$  by  $\hat{J}(i, z)$

# BELLMAN EQN METHODS

Idea : Solve Directly BELLMAN EQN By USING A NN.

$$\min_{\pi} \sum_{i \in S} (\tilde{J}(i, \pi) - \min_u \sum_{j} \pi_{ij}(u) (g_{ij}(u) + \tilde{J}(j, \pi)))^2$$

RESIDUAL OF BELLMAN EQN. ( $J^* = T J^*$ )

RESIDUAL :  $D(i, \pi) = \tilde{J}(i, \pi) - \min_u \sum_j \pi_{ij}(u) (\underline{g_{ij}(u)} + \tilde{J}(j, \pi))$

$\Leftrightarrow \min_{\pi} \sum_{i \in S} D(i, \pi)^2$

APPLYING SGD TO MINIMIZE THE RESIDUAL

$$\begin{aligned}\pi_{k+1} &= \pi_k - \gamma \underbrace{\nabla G(\pi_k)}_{\text{? SAMPLE AT } k} \nabla D(i, \pi_k) \\ &= \pi_k - \gamma \nabla D(i, \pi_k) \left( \sum_j p_{ij}(u) \underbrace{\nabla_j g}_{\text{?}}(j, \pi_k) - \underbrace{\nabla_i g}_{\text{?}}(i, \pi_k) \right)\end{aligned}$$

WHERE

$$\overline{u} = \arg \min_u \sum_j p_{ij}(u) (g(i, u, j) + \overbrace{J(g, \pi_k)})$$







